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Approach for Sensing Coverage Using Border Effect in Wireless Sensor Network

Geeta Rani and Nasib Singh Gill
Department of Computer Science & Applications Maharshi Dayanand University Rohtak-124001 Haryana, India
Corresponding Author: Geeta Rani

ABSTRACT**:** *It has been observed that coverage is a basic issue in remote sensor organize (WSN).This issue can be elucidated by survey the less number of the handle vivacious at the specific time. The sensors from the dynamic groups defines all objectives and at danger to transmit the information to the base station and remaining center points are in rest mode. This paper discusses a couple of critical computations - Modified K-Means, PSO figuring and AOMDV utilized as a bit of remote sensor systems. The Modified K-Means is utilized to scope the range, AOMDV used for multipath coordinating and PSO is utilized to locate the perfect way.*

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I. Introduction

A wireless sensor network (WSN) is a wireless network comprising of spatially circulated self-governing gadgets utilizing sensors to screen physical or natural conditions. A WSN framework joins a passage that gives wireless network back to the wired world and dispersed hubs.

Late years have seen tremendous movement in remote sensor systems due to diminishment being created costs and suddenness in gear manufacturing. Past a couple of decades have been separate with energetic usage of remote sensor arranges in a gathering of fields. By and by remote sensor systems are used as a piece of military shadowing, an area checking, seismic activity perception and in indoor applications. These remote sensors have given us gadget to screen a district of interest remotely. Each one of the one ought to do send these sensors, aerodynamically or physically, and a while later these sensors which outline the crossing points of the system amass information from the domain under investigation. The information along these lines got is trade back to the "central server" or "base station" where the information is taken care of. The base server is every so often connected with Internet which by then exchanges the dealt with information by methods for satellite to the essential station or control group for moreover planning and examination. No or no dealing with is done while data is traded from convergences.

Sensor crossing points that make the remote system are self-decision convergences with a microcontroller, no less than one sensor, a handset, actuators and a battery for control supply. These crossing points, as a rule furthermore suggested bits. These sensors have no memory and perform little measure of taking care of with the data procured. Directly isolated from checking, assembling and transmitting data beginning with one convergence then onto the following and to the base station additionally the taking care of unit oversees and controls handiness of various portions of the sensor crossing point.. By the by the memory operation is an overhead too. This is by virtue of the sensors are outfitted with a battery which much of the time is non-replaceable. Along these lines augment in getting ready would saw greater essentialness is being eaten up and from now on sensor lifetime would reduce thusly impacting the lifetime of the system. As determined before the exchanging of data is done by following a particular transmission tradition. However this office is expert by the transmission unit of the sensor. Commonly the sensor has a handset that can go about as both transmitter and an authority. The transmitter and the recipient gear both are not kept separate with a particular true objective to save space and essentialness. These days, sensors can give through transmission media reaching out from tremendous electromagnetic range. A remote sensor organize is to be found in one of the two ways: arranged and unintended. In the organized technique for sending a specific number of sensors are placed in imperative concentrations in fated way.

Here it should be seen that the region to be checked can be gotten to physically, thusly the cost isn't a factor under such conditions. These crossing points are put using a fated count with the true objective that the district to be secured is increased setting less overhead on transmission and battery thusly growing the system lifetime. The remote sensor arrange confront distinctive issues one of which fuses extent of the given district under obliged essentialness. This issue of increasing the system lifetime while following the extension and imperativeness parameters or obstruction is known as the Target Coverage Problem in Wireless Sensor Networks.

As the sensor crossing points are battery driven so they have obliged essentialness also and thus the essential test winds up discernibly extending the degree locale and besides guaranteeing a drawn out system lifetime. The work has been done to address this issue yet basically as the trial obviously contain time constraint; therefore the issue advances toward getting to be time subordinate, which hence is non-polynomial in nature. By and by even non-coordinate issues have a place with the NP-Hard class thus only two or three heuristics have been prescribed to address the Target Coverage Problem if not in perfect, by then close perfect or dangerous time. One of such computation is inspects in also range which is used as a gage against our proposed count.

1.1 MODIFIED K-MEANS:

This paper presents a data bundling approach using balanced K-Means computation in perspective of the difference in the affectability of starting concentration of gatherings. This count fragments the whole space into different areas and determines the repeat of data point in each part. The section which shows most noteworthy repeat of data point will have the most extraordinary probability to contain the centroid of gathering. The amount of pack's centroid (k) will be given by the customer comparatively like the traditional K-mean figuring and the amount of division will be k*k (`k' vertically and moreover `k' on a level plane). In case the most critical repeat of data point is same in different parts and the upper bound of bit crosses the farthest point `k' by then focalizing of different segments wind up detectably required and a while later take the most astonishing k parcel for discovering the basic centroid of packs. In this paper we moreover portray an edge expel for each bundle's centroid to examine the division between data point and gathering's centroid with this cutoff evacuate through which we can restrain the computational effort in the midst of estimation of partition between data point and gathering's centroid. It is exhibited that how the changed k-mean figuring will lessen the multifaceted nature and the effort of numerical calculation, keeping up the simplicity of executing the k-mean computation. It allocates the data point to their appropriate class or bundle more effectively.

We have presented a changed k-means computation which takes out the issue of period of fumes gatherings (with a couple of uncommon cases). Here, the fundamental structure of the principal k-suggests is secured nearby all its critical properties. Another inside vector computation framework engages us to rename the bundling strategy and to accomplish our target. The balanced count is found to work alluringly, with some unforeseen exceptional cases which are to a great degree unprecedented for all intents and purposes.

Modified approach K-mean calculation:

The K-mean estimation is a pervasive packing count and has its application in data mining, picture division, bioinformatics and various distinctive fields. This figuring capacities outstandingly with little datasets. In this paper we proposed a count that capacities outstandingly with broad datasets. Changed k-mean count declines getting into locally perfect course of action in some degree, and decreases the assignment of gathering - botch measure.

Algorithm: Modified approach (S, k), S={x1,x2,…,xn }

Input: The number of clusters k1 (k1> k) and a dataset containing n objects(Xij+).

Output: A set of k clusters (Cij) that minimize the Cluster - error criterion.

Algorithm

1. Compute the distance between each data point and all other data- points in the set D
2. Find the closest pair of data points from the set D and form a data-point set Am (1<= p <= k+1) which contains these two data- points, Delete these two data points from the set D
3. Find the data point in D that is closest to the data point set Ap, Add it to Ap and delete it from D
4. Repeat step 4 until the number of data points in Am reaches (n/k)
5. If p<k+1, then p = p+1, find another pair of data points from D between which the distance is the shortest, form another data-point set Ap and delete them from D, Go to step 4.

1.2 PSO Algorithm

Particle swarm advancement was exhibited by Kennedy and Eberhart (1995). It has builds up in the entertainment of social works on using instruments and musings taken from PC plans and social cerebrum inquire about explore.

The target of the count is to have each one of the particles discover the optima in a multi-dimensional hyper-volume. This is refined by apportioning at first unpredictable positions to all particles in the space and minimal basic self-assertive velocities. The computation is executed like a reenactment, pushing the position of each molecule in this manner in perspective of its speed, the best known overall position in the issue space and the best position known to a molecule. The objective work is assessed after each position revive. After some time, through a mix of examination and abuse of known incredible positions in the interest space, the particles group or get together around optima, or a couple of optima.

1.2.1 Algorithm Outline

The particle swarm algorithm begins by creating the initial particles, and assigning them initial velocities.

- It evaluates the objective function at each particle location, and determines the best (lowest) function value and the best location.
- It chooses new velocities, based on the current velocity, the particles' individual best locations, and the best locations of their neighbors.
- It then iteratively updates the particle locations (the new location is the old one plus the velocity, modified to keep particles within bounds), velocities, and neighbors.
- Iterations proceed until the algorithm reaches a stopping criterion.

1.2.2 The algorithm

As communicated some time as of late, PSO imitates the acts of winged creature running. Accept the going with circumstance: a get-together of feathered animals are self-assertively looking food in a range. There is only a solitary piece of food in the domain being looked. Each one of the winged animals don't know where the sustenance is. In any case, they know how far the sustenance is in each cycle. So what's the best method to find the food? The capable one is to take after the feathered animal which is nearest to the sustenance.

PSO picked up from the circumstance and used it to deal with the headway issues. In PSO, each single course of action is a "feathered animal" in the interest space. We call it "atom". All of particles have wellbeing regards which are evaluated by the health ability to be improved, and have speeds which facilitate the flying of the particles. The particles fly through the issue space by following the present perfect particles.

PSO is presented with a social event of sporadic particles (courses of action) and after that searches for optima by invigorating times. In each cycle, each atom is invigorated by following two "best" qualities. The first is the best course of action (wellbeing) it has achieved up until this point. (The health regard is moreover secured.) This regard is called pbest. Another "best" regard that is trailed by the atom swarm enhancer is the best regard, got so far by any particle in the people. This best regard is an overall best and called gbest. Exactly when a particle expels a bit of the masses as its topological neighbors, the best regard is an adjacent best and is called lbest.

The pseudo code of the PSO is as follows

```
for each particle
   Initialize particle
END
Do
   For each particle
      Calculate fitness value
      If the fitness value is better than the best fitness value (pBest) in history
set current value as the new pBest
   End
   Choose the particle with the best fitness value of all the particles as the gBest
   For each particle
      Calculate particle velocity according equation (a)
      Update particle position according equation (b)
   End
While maximum iterations or minimum error criteria is not attained
```

Particles' speeds on each measurement are clasped to a most extreme speed Vmax. On the off chance that the aggregate of increasing speeds would make the speed on that measurement surpass Vmax, which is a parameter indicated by the client. At that point the speed on that measurement is constrained to Vmax.

II. Related Work

Aneesh Kumar V.N., AyanaAjith [1], A wireless sensor network is commonly used to monitoring and recording special events in a geographical area with the help of number of sensors called sensor nodes. These sensor nodes are small in size, weight and portability. They are very vulnerable to various type of failures. These failures form holes in the coverage area. The four key elements that ensure coverage for WSNs are determining the boundary of RoI, detecting coverage holes and estimating their characteristics, determining the best target locations to relocate mobile nodes to repair holes, and dispatching mobile nodes to the target location while minimizing the moving and messaging cost. The coverage enhancement and hole healing is a big task in the field of wireless sensor networks. There are different methods are available for detecting holes and their boundary. Also different methods are used to enhance coverage area and whole healing. These works goes through the available methods for this purpose and differentiate their performance.

Feng Li, Jun Luo [2], Although the problem of k-area coverage has been intensively investigated for dense wireless sensor networks (WSNs), how to arrive at a k-coverage sensor deployment that optimizes certain objectives in relatively sparse WSNs still faces both theoretical and practical difficulties. In this paper, we present a practical algorithm LAACAD (Load balancing k-Area Coverage through Autonomous Deployment) to move sensor nodes toward k-area coverage, aiming at minimizing the maximum sensing range required by the nodes. LAACAD enables purely autonomous node deployment as it only entails localized computations. We prove the convergence of the algorithm, as well as the (local) optimality of the output. We also show that our optimization objective is closely related to other frequently considered objectives. Therefore, our practical algorithm design also contributes to the theoretical understanding of the k-area coverage problem. Finally, we use extensive simulation results both to confirm our theoretical claims and to demonstrate the efficacy of LAACAD.

A. Capone, M. Cesana, D. De Donno, I. Filippini [3], Data collected by sensors often have to be remotely delivered through multi-hop wireless paths to data sinks connected to application servers for information processing. The position of these sinks has a huge impact on the quality of the specific Wireless Sensor Network (WSN). Indeed, it may create artificial traffic bottlenecks which affect the energy efficiency and the WSN lifetime. This paper considers a heterogeneous network scenario where wireless sensors deliver data to intermediate gateways geared with a diverse wireless technology and interconnected together and to the sink. An optimization framework based on Integer Linear Programming (ILP) is developed to locate wireless gateways minimizing the overall installation cost and the energy consumption in the WSN, while accounting for multi-hop coverage between sensors and gateways, and connectivity among wireless gateways. The proposed ILP formulations are solved to optimality for medium-size instances to analyze the quality of the designed networks, and heuristic algorithms are also proposed to tackle large-scale heterogeneous scenarios.

KoenLangendoen ,NielsReijers [4], This paper studies the problem of determining the node locations in ad-hoc sensor networks. We compare three distributed localization algorithms (Ad-hoc positioning, Robust positioning, and N-hop multilateration) on a single simulation platform. The algorithms share a common, three-phase structure: (1) determine node–anchor distances, (2) compute node positions, and (3) optionally refine the positions through an iterative procedure. We present a detailed analysis comparing the various alternatives for each phase, as well as a head-to-head comparison of the complete algorithms. The main conclusion is that no single algorithm performs best; which algorithm is to be preferred depends on the conditions (range errors, connectivity, anchor fraction, etc.). In each case, however, there is significant room for improving accuracy and/or increasing coverage.

Xinbing Wang, Sihui Han, Yibo Wu, and Xiao Wang [5],In this paper, we investigate the coverage and energy consumption control in mobile heterogeneous wireless sensor networks (WSNs). By term heterogeneous, we mean that sensors in the network have various sensing radius, which is an inherent property of many applied WSNs. Two sensor deployment schemes are considered –uniform and Poisson schemes. We study the asymptotic coverage under uniform deployment scheme with i.i.d. and 1-dimensional random walk mobility model, respectively. We propose the equivalent sensing radius (ESR) for both cases and derive the critical ESR correspondingly. Our results show that the network performance largely depends on ESR. By controlling ESR, we can always promise the network achieve full coverage, regardless of the total number of sensors or the sensing radius of a single senor under random mobility patterns, which is a much easier and more general way to operate coverage control. Meanwhile, we can operate a tradeoff control between coverage performance and energy consumption by adjusting ESR. We demonstrate that 1-dimensional random walk mobility can decrease the sensing energy consumption under certain delay tolerance, though requires larger ESR. Also, we characterize the role of heterogeneity in coverage and energy performance of WSNs under these two mobility models, and present the discrepancy of the impact of heterogeneity under different models. Under the Poisson deployment scheme, we investigate dynamic k-coverage of WSNs with 2-dimensional random walk mobility model. We present the relation between network coverage and the sensing range, which indicates how coverage varies according to sensing capability. Both k-coverage at an instant and over a time interval are explored and we derive the expectation of fraction of the whole operational region that is k-covered, which also identifies the coverage improvement brought by mobility.

Isabel Dietrich, Falko Dressler [6], Network lifetime has become the key characteristic for evaluating sensor networks in an application specific way. Especially the availability of nodes, the sensor coverage, and the connectivity have been included in discussions on network lifetime. Even quality of service measures can be reduced to lifetime considerations. A great number of algorithms and methods were proposed to increase the lifetime of a sensor network { while their evaluations were always based on a particular definition of network lifetime. Motivated by the great differences in existing definitions of sensor network lifetime that are used in relevant publications, we reviewed the state of the art in lifetime definitions, their differences, advantages, and limitations. This survey was the starting point for our work towards a generic definition of sensor network lifetime for use in analytic evaluations as well as in simulation models { focusing on a formal and concise definition of accumulated network lifetime and total network lifetime. Our definition incorporates the components of existing lifetime definitions, and introduces some additional measures. One new concept is the ability to express the service disruption tolerance of a network. Another new concept is the notion of time-integration: in many cases, it is sufficient if a requirement is fulfilled over a certain period of time, instead of at every point in time. In addition, we combine coverage and connectivity to form a single requirement called connected coverage. We show that connected coverage is different from requiring non-combined coverage and connectivity. Finally, our definition also supports the concept of graceful degradation by providing means of estimating the degree of compliance with the application requirements. We demonstrate the applicability of our definition based on the surveyed lifetime definitions as well as using some example scenarios to explain the various aspects influencing sensor network lifetime.

LoukasLazos and RadhaPoovendran [7], We study the problem of coverage in planar heterogeneous sensor networks. Coverage is a performance metric that quantifies how well a field of interest is monitored by the sensor deployment. To derive analytical expressions of coverage for heterogeneous sensor networks, we formulate the coverage problem as a set intersection problem, a problem studied in integral geometry. Compared to previous analytical results, our formulation allows us to consider a network model where sensors are deployed according to an arbitrary stochastic distribution; sensing areas of sensors need not follow the unit disk model but can have any arbitrary shape; sensors need not have an identical sensing capability. Furthermore, our formulation does not assume deployment of sensors over an infinite plane and, hence, our derivations do not suffer from the border effect problem arising in a bounded field of interest. We compare our theoretical results with the spatial Poisson approximation that is widely used in modeling coverage. By computing the Kullback-Leibler and total variation distance between the probability density functions derived via our theoretical results, the Poisson approximation, and the simulation, we show that our formulas provide a more accurate representation of the coverage in sensor networks. Finally, we provide examples of calculating network parameters such as the network size and sensing range in order to achieve a desired degree of coverage.

Mohamed YounisKemalAkkaya [8], The major challenge in designing wireless sensor networks (WSNs) is the support of the functional, such as data latency, and the non-functional, such as data integrity, requirements while coping with the computation, energy and communication constraints. Careful node placement can be a very effective optimization means for achieving the desired design goals. In this paper, we report on the current state of the research on optimized node placement in WSNs. We highlight the issues, identify the various objectives and enumerate the different models and formulations. We categorize the placement strategies into static and dynamic depending on whether the optimization is performed at the time of deployment or while the network is operational, respectively. We further classify the published techniques based on the role that the node plays in the network and the primary performance objective considered. The paper also highlights open problems in this area of research.

Jennifer C. Hou, David K. Y. Yau [9], Ad-hoc networks of devices and sensors with (limited) sensing and wireless communication capabilities are becoming increasingly available for commercial and military applications. The first step in deploying these wireless sensor networks is to determine, with respect to application-specific performance criteria, (i) in the case that the sensors are static, where to deploy or activate them; and (ii) in the case that (a subset of) the sensors are mobile, how to plan the trajectory of the mobile sensors. These two cases are collectively termed as the coverage problem in wireless sensor networks. In this book chapter, we give a comprehensive treatment of the coverage problem. Specifically, we first introduce several fundamental properties of coverage that have been derived in the literature and the corresponding algorithms that will realize these properties. While giving insights on how optimal operations can be devised, most of the properties are derived (and hence their corresponding algorithms are constructed) under the perfect disk assumption. Hence, we consider in the second part of the book chapter coverage in a more realistic setting, and allow (i) the sensing area of a sensor to be anisotropic and of arbitrary shape, depending on the terrain and the meteorological conditions, and (ii) the utilities of coverage in different parts of the monitoring area to be non-uniform, in order to account for the impact of a threat on the population, or the likelihood of a threat taking place at certain locations. Finally, in the third part of the book chapter, we consider mobile sensor coverage, and study how mobile sensors may navigate in a deployment area in order to maximize threat-based coverage.

Ji Li, Lachlan L.H. Andrew, ChuanHengFoh, Moshe Zukerman [10], Wireless communication between sensors allows the formation of flexible sensor networks, which can be deployed rapidly over wide or inaccessible areas. However, the need to gather data from all sensors in the network imposes constraints on the distances between sensors. This survey describes the state of the art in techniques for determining the minimum density and optimal locations of relay nodes and ordinary sensors to ensure connectivity, subject to various degrees of uncertainty in the locations of the nodes.

Mohamed Hefeeda and HosseinAhmadi [11], We propose a new probabilistic coverage protocol (denoted by PCP) that considers probabilistic sensing models. We design PCP keeping in mind that no single sensing model (probabilistic or not) will accurately model all types of sensors in all environments. It is expected that different sensor types will require different sensing models. Even for the same sensor type, the sensing model may need to be changed in different environments. Designing, implementing, and testing a different coverage protocol for each sensing model is indeed an extremely costly process, if at all possible. To address this challenging task, we design our protocol with limited dependence on the sensing model. In particular, our protocol requires the computation of a single parameter from the adopted sensing model, while everything else remains the same. We show how this parameter can be derived in general, and we actually do the calculations for two example sensing models: (i) the probabilistic exponential sensing model, and (ii) the commonly-used deterministic disk sensing model. The first model is chosen because it is conservative in terms of estimating sensing capacity, and it has been used before in another probabilistic coverage protocol, which enables us to conduct a fair comparison. Because it is conservative, the exponential sensing model can be used as a first approximation for many other sensing models. The second model is chosen to show that our protocol can easily function as a deterministic coverage protocol. In this case, we compare our protocol against two recent deterministic protocols that were shown to outperform others in the literature. Our comparisons indicate that our protocol outperforms all other protocols in several aspects, including number of activated sensors and total energy consumed. We also demonstrate the robustness of our protocol against random node failures, node location inaccuracy, and imperfect time synchronization.

Ashraf Hossain, P. K. Biswas, S. Chakrabarti [12], Network coverage of wireless sensor network (WSN) means how well an area of interest is being monitored by the deployed network. It depends mainly on sensing model of nodes. In this paper, we present three types of sensing models viz. Boolean sensing model, shadow-fading sensing model and Elfes sensing model. We investigate the impact of sensing models on network coverage. We also investigate network coverage based on Poisson node distribution. A comparative study between regular and random node placement has also been presented in this paper. This study will be useful for coverage analysis of WSN.

III. Proposed Methodology

We proposed the new procedure which tackled every one of the issues portray in issue explanation. Primary destinations of our proposed framework are ideal way, target scope, increment the life time of the system and disappointment recognition. We present the novel approach in which we combine the different systems called AOMDV, Modified K-Means and PSO (Particle Swarm Optimization) to discover Optimum way, target scope and disappointment location in remote sensor Network. We have proposed a System which depends on upon mix approach of utilizing ideal way using adjusted k-implies for territory scope and AOMDV for multipath directing and afterword applying PSO for briefest way to discover best outcome.

Step 1: The primary stage applies the K-implies calculation to segment the system into k groups.

Stage 2: Next, the AOMDV calculation for the multipath directing inside each group got by the Modified K-Means.

Stage 3: Finally, the PSO calculation looks for the best CH (group head) inside each bunch and assesses the ideal way.

With our proposed calculation, the grouping issue will be effortlessly made do with less calculation required and we can discover ideal way, target scope, lessen vitality utilization and increment organize life time with disappointment discovery.

IV. Result Analysis/Implementation:

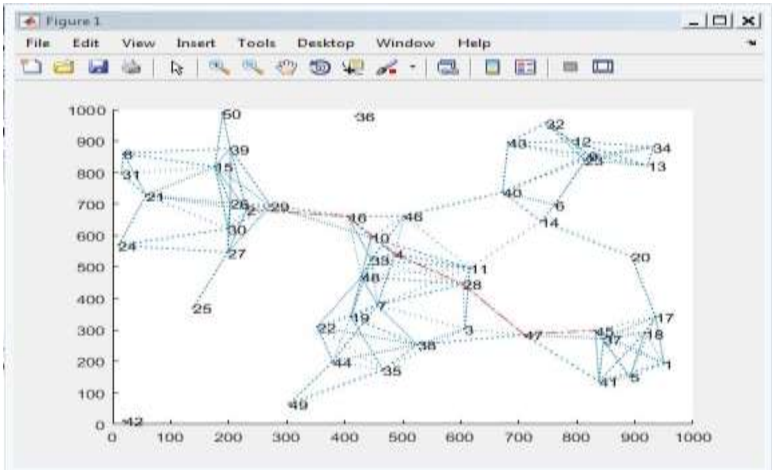

Fig 4.1: Optimize Path between source & destination node

Figure 4.1 used for displaying UI of path optimization process. In this figure node 2 is given as source node and node 45 as destination node. The red line indicates the optimized path between given source node and destination node.

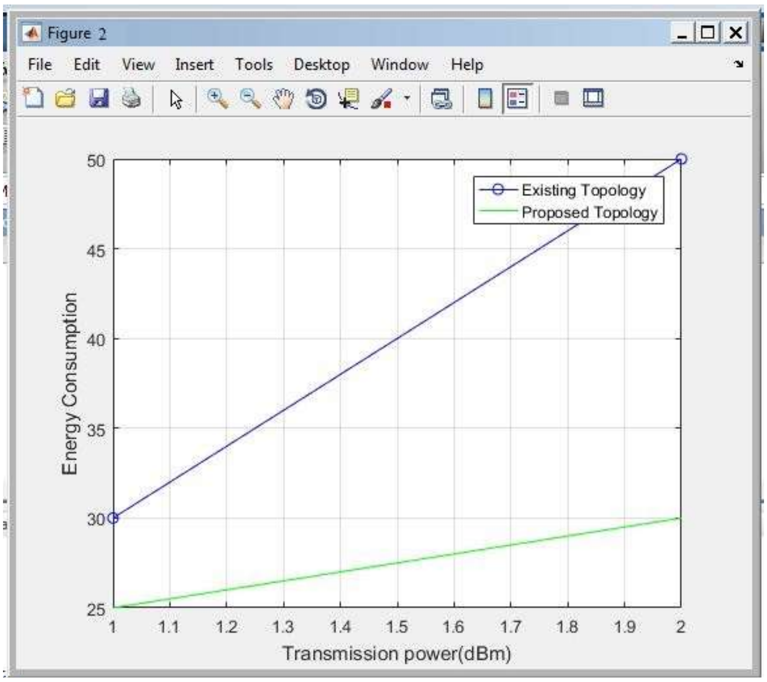

Figure 4.2: Comparison graph between existing and proposed topology according to energy & transmission power

```
path =

     2    16     4    28    47    45


cost =

     5
```

Figure 4.3: Final cost

V. Conclusion

Ensuing to assessing some in advance done basic investigates we have proposed a theoretical working framework which depends upon mix approach of utilizing perfect way using modified k-infers for an area scope and a brief timeframe later apply PSO for briefest system course close by to discover culminate result. Along these lines the paper finishes a structure that can illuminate the issues of WSN which are degree issue, perfect way, and essentialness capable system.

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